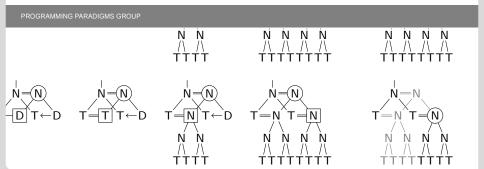


dup – Explicit un-sharing in Haskell

Haskell Implementors Workshop 2012 - Lightning Talk



My motivation



We need to provide our programmers with better tools to

analyze

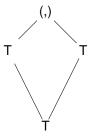
and

control

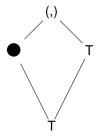
the space behaviour of their Haskell programs.



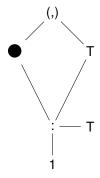




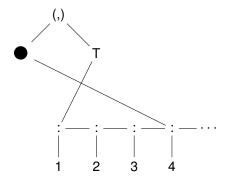




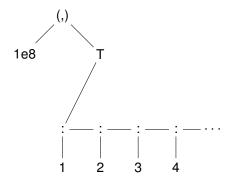






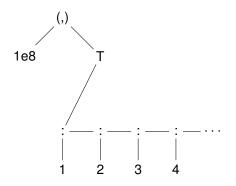








```
let xs = [1..100000000]
in (last xs, length xs)
```



the programmer might want to avoid to have the list shared



```
let xs () = [1..100000000]
in (last $ xs (), length $ xs ())
```



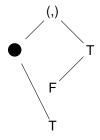


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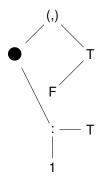


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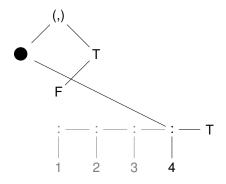


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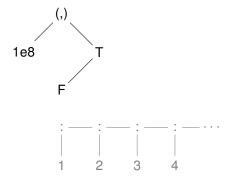


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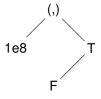


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works, but fragile - might be thwarted by compiler optimizations

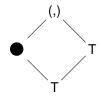


```
let xs = [1..1000000000]
in (case dup xs of Box xs' -> last xs',
    case dup xs of Box xs' -> length xs')
```



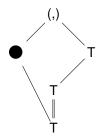


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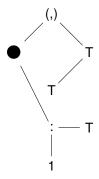


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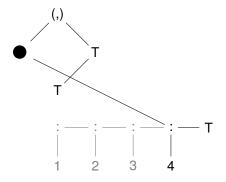


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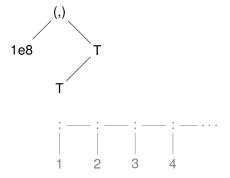


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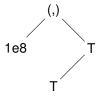


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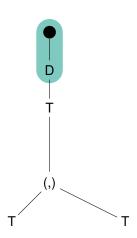


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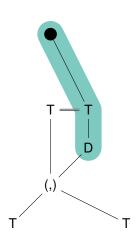


the consumer, not the generator, controls sharing. no code restructuring.

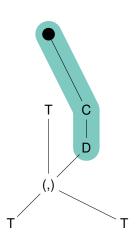




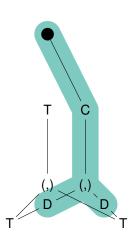




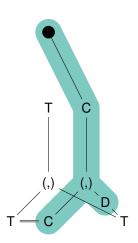




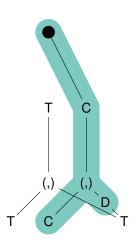




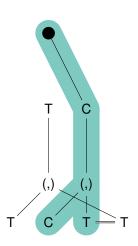




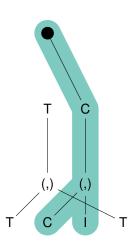






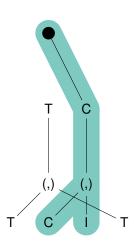








morally, deepDup x copies the whole heap reachable by x



Comes with proofs included.



$$\frac{\Gamma, x \mapsto e, x' \mapsto \hat{e} : x' \Downarrow \Delta : z \qquad x' \text{ fresh}}{\Gamma, x \mapsto e : \operatorname{dup} x \Downarrow \Delta : z} \operatorname{Dup}$$

$$\begin{array}{c} x' \mapsto \hat{\mathbf{e}}[y_1'/y_1, \ldots, y_n'/y_n], \\ \Gamma, x \mapsto e, \quad y_1' \mapsto \operatorname{deepDup} y_1, \ldots, y_n' \mapsto \operatorname{deepDup} y_n : x' \Downarrow \Delta : z \\ \hline \frac{\operatorname{ufv}(e) = \{y_1, \ldots, y_n\} \qquad x', \ y_1', \ldots, y_n \operatorname{fresh}}{\Gamma, x \mapsto e : \operatorname{deepDup} x \Downarrow \Delta : z} \\ \end{array} \text{Deep}$$

(based on Launchbury's "A natural seantics for lazy evaluation")

Where to read more



See

http://arxiv.org/abs/1207.2017

for

- more motiviation,
- benchmarked comparison with other approaches to avoid sharing,
- semantics and proofs,
- details on the implementation and
- a description of current shortcomings.

See

http://darcs.nomeata.de/ghc-dup

for

the code.

A related, younger idea



```
import GHC.Prim (noupdate)
```

```
let xs = noupdate [1..100000000]
in (last xs, length xs)
```

For a thunk wrapped in noupdate :: $a \rightarrow a$, no blackhole and no update frame is created ⇒ sharing is effectively prevented.

(Ask me for my ghc branch.)

Also nice: ghc-vis



Demonstration

see

http://hackage.haskell.org/package/ghc-vis

and

http://felsin9.de/nnis/ghc-vis/